

Amendments to the Claims

This listing of claims will replace all prior versions, and listings, of claims in the application:

1. (currently amended) A system for organizing and accessing a database, the system

comprising:

a processor operable to execute computer program instructions;

a memory operable to store the computer program instructions executable by the
processor; and

the computer program instructions stored in the memory and executable to
perform the steps of:

providing a secondary B+tree for indexing a primary B+tree, wherein the secondary B+tree
comprises a plurality of mapping table row identifiers, where each mapping row identifier is an entry
in a row of the secondary B+tree that uniquely identifies a row within a mapping table and a first
plurality of database addresses for leaf blocks of the primary B+tree, and

creating the mapping table with a corresponding number of database addresses for leaf blocks
of the primary B+tree, wherein each of the database addresses is a ~~rows each comprising an index~~
~~key value, and~~ a guess-database address that is a database address , which corresponds to a guess as
to what address block a row of the primary B+tree a row may be found, where data stored in the row
~~of the primary B+ tree of the database is retrieved using the secondary B+tree index.~~

2. (previously presented) The system according to claim 1, wherein the guess-database address
are 4 bytes of address blocks in the primary B+tree.

3. (cancelled)

4. (previously presented) The system according to claim 1, further comprising:

a guess-database address quality statistic for the secondary index, where the guess-database address quality statistic represents a ratio of how often the guesses as to where rows may be found in an address block of the primary B+tree are accurate.

5. (currently amended) A computer-implemented method for managing a database system, the method comprising:

creating a secondary B+tree for indexing a primary B+tree, wherein the secondary B+tree comprises a plurality of mapping table row identifiers, where each mapping row identifier is an entry in a row of the secondary B+tree that uniquely identifies a row within a mapping table and a first plurality of database addresses for leaf blocks of the primary B+tree, and

creating the mapping table with a corresponding number of database addresses for leaf blocks of the primary B+tree, wherein each of the database addresses is a ~~rows each comprising an index key value, and~~ a guess-database address that is a database address, which corresponds to a guess as to what address block a row of the primary B+tree ~~a row may be found, where data stored in the row of the primary B+ tree of the database is retrieved using the secondary B+tree index.~~

6. (previously presented) The method according to claim 18, further comprising:

inserting a row of the secondary index, wherein inserting the row comprises inserting a row comprising an index key value, a mapping table rowid value and a guess database address.

7. (previously presented) The method according to claim 18, further comprising:

deleting a row of the secondary index, wherein deleting the row comprises locating a row comprising an index key value and a mapping table rowid value and deleting the row.

8. (previously presented) The method according to claim 18, further comprising:

updating the secondary index, wherein updating the secondary index comprises locating a row of the secondary index comprising an old index key value and a mapping table rowid value,

deleting the row and inserting in the row a new index key value, a mapping table rowid value and a guess database address.

9. (currently amended) The method according to claim 18, wherein retrieving data stored in the database system further comprises:

obtaining a first guess database address ~~value~~ representing a first address block of the primary B+tree structure;

searching the first address block of the primary B+tree for a row that contains a mapping table rowid value that is the same as a mapping table rowid value in the row where the first guess database address is stored in the secondary index row; and

if the mapping table rowid is found then the correct row in the primary B+tree has been located and the data is retrieved.

10. (previously presented) The method according to claim 9, wherein retrieving data stored in the database system further comprises:

if the mapping table rowid value is not found, then utilizing the mapping table rowid value stored in the row where the first guess database address is stored in the secondary index to access a mapping table row stored in a mapping table;

utilizing a second guess-database address stored in the mapping table row to access a second address block of the primary B+tree;

searching the second address block of the primary B+tree for a primary key that matches a primary key stored in the mapping table row; and

if the primary key is found, then the data is retrieved.

11. (previously presented) The method according to claim 10, wherein if the primary key is not found retrieving data further comprises:

traversing the primary B+tree utilizing the primary key value from the mapping table row to identify the database address to complete the query.

12. (previously presented) The method according to claim 11, further comprising:

maintaining a guess-database address quality statistic for the secondary index;

maintaining a guess-database address quality statistic for the mapping table;

utilizing the statistics to assess guess-database address quality; and

carrying out the query based upon guess-database quality in the secondary index and mapping table, where each of the guess-database address quality statistic represents a ratio of how often the guesses as to where rows may be found in an address block of the primary B+tree are accurate.

13. (original) The method according to claim 12, further comprising:

estimating guess-database address quality;

estimating the cost of the query based upon the estimated guess-database address quality; and

carrying out the query starting with an index structure with the lowest estimated cost.

14. (previously presented) The method according to claim 5, wherein the guess-database address are 4 bytes of address blocks in the primary B+tree.

15. (currently amended) A computer program product for performing a process of managing a database system, comprising:

a computer readable storage medium; and

computer program instructions, recorded on the computer readable storage medium,

executable by a processor, for performing the steps of:

creating a secondary B+tree for indexing a primary B+tree, wherein the secondary B+tree comprises a plurality of mapping table row identifiers, where each mapping row identifier is an entry in a row of the secondary B+tree that uniquely identifies a row within a mapping table and a first plurality of database addresses for leaf blocks of the primary B+tree, and

creating the mapping table with a corresponding number of database addresses for leaf blocks of the primary B+tree, wherein each of the database addresses is a ~~rows each comprising an index key value, and~~ a guess-database address that is a database address , which corresponds to a guess as to what address block a row of the primary B+tree ~~a row may be found, where data stored in the row of the primary B+ tree of the database is retrieved using the secondary B+tree index.~~

16. (currently amended) A system for performing a database management process, comprising

a processor operable to execute computer program instructions; and

a memory operable to store computer program instructions executable by the processor,

for performing the steps of:

creating secondary B+tree for indexing a primary B+tree, wherein the secondary B+tree comprises a plurality of mapping table row identifiers, where each mapping row identifier is an entry in a row of the secondary B+tree that uniquely identifies a row within a mapping table and a first plurality of database addresses for leaf blocks of the primary B+tree, and

creating the mapping table with a corresponding number of database addresses for leaf blocks of the primary B+tree, wherein each of the database addresses is a~~rows each comprising an index key value, and~~ a guess-database address that is a database address , which corresponds to a guess as to what address block a row of the primary B+tree ~~a row~~ may be found,~~where data stored in the row of the primary B+ tree of the database is retrieved using the secondary B+tree index.~~

17. (previously presented) The ~~system~~ method according to claim 1, wherein each row in the plurality of rows further comprising a mapping table rowid value that identifies a row within a mapping table

18. (previously presented) The method according to claim 5, wherein each row in the plurality of rows further comprising a mapping table rowid value that identifies a row within a mapping table.

19. (previously presented) The computer program product according to claim 15, wherein each row in the plurality of rows further comprising a mapping table rowid value that identifies a row within a mapping table.

20. (previously presented) The system according to claim 16, wherein each row in the plurality of rows further comprising a mapping table rowid value that identifies a row within a mapping table.